HLIBCov: Parallel Hierarchical Matrix Approximation of Large Covariance Matrices and Likelihoods with Applications in Parameter Identification

Alexander Litvinenko

King Abdullah University of Science and Technology, Thuwal 23955-6900, Saudi Arabia

Abstract. The main goal of this article is to introduce the parallel hierarchical matrix library HLIBpro to the statistical community. We describe the HLIBCov package, which is an extension of the HLIBpro library for approximating large covariance matrices and maximizing likelihood functions. We show that an approximate Cholesky factorization of a dense matrix of size $2M \times 2M$ can be computed on a modern multi-core desktop in few minutes. Further, HLIBCov is used for estimating the unknown parameters such as the covariance length, variance and smoothness parameter of a Matérn covariance function by maximizing the joint Gaussian log-likelihood function. The computational bottleneck here is expensive linear algebra arithmetics due to large and dense covariance matrices. Therefore covariance matrices are approximated in the hierarchical ($H$-) matrix format with computational cost $O(k^2n\log^2 n/p)$ and storage $O(kn\log n)$, where the rank $k$ is a small integer (typically $k < 25$), $p$ the number of cores and $n$ the number of locations on a fairly general mesh. We demonstrate a synthetic example, where the true values of known parameters are known. For reproducibility we provide the C++ code, the documentation, and the synthetic data.

Key words: Computational statistics; parallel hierarchical matrices; large datasets; Matérn covariance; random fields; spatial statistics, HLIB, HLIBpro, Cholesky, matrix determinant

HLIBCov software library

Program title: HLIBCov

http://www.global-sci.com/
Nature of problem: To estimate the unknown parameters (variance, smoothness parameter, and covariance length) of a covariance function by maximizing the joint Gaussian log-likelihood function with a log-linear computational cost and storage.

Software licence: GPL 2.0

CiCP scientific software URL:

Distribution format: *.cc files via github

Programming language(s): C++

Computer platform: 32 or 64

Operating system: Unix-like operating systems

Compilers: clang

RAM: 4GB and more (depending on the matrix size)

External routines/libraries: HLIBpro requires:

- BLAS/LAPACK,
- Intel’s Threading Building Blocks [http://www.threadingbuildingblocks.org/],
- Boost [http://www.boost.org],
- zlib [http://zlib.net],
- FFTW3 (system dependent, [http://www.fftw.org]),
- Scotch (system dependent, [http://gforge.inria.fr/projects/scotch/]),
- METIS (system dependent, [http://glaros.dtc.umn.edu/gkhome/views/metis])
- SCons [http://www.scons.org]

HLIBCov requires: GNU Scientific Library [https://www.gnu.org/software/gsl/]

Running time: $O(k^2n\log^2 n)/p$ with $p$ number of cores

Restrictions: 1. For multiple-core architectures with shared memory systems
2. Compiling with GCC 5 will generate linking errors.

Supplementary material and references: [www.HLIBpro.com] and references therein.

Additional Comments: HLIBpro is a software library implementing parallel algorithms for Hierarchical matrices. It is freely available in binary form for academic purposes. HLIBpro algorithms design for 1, 2 and 3-dimensional problems.
1 Introduction

This paper is complementary to the paper [2].

**Novelty of this work.** In this paper we use parallel hierarchical (H-) matrices for approximating dense covariance matrices and the joint Gaussian log-likelihood with computational complexity $O(k^\alpha n \log^\alpha n / p)$, where $p$ is the number of cores, $n$ is the number of measurements; $k \ll n$ is the maximal rank, used in the hierarchical matrix, which defines the quality of the approximation; and $\alpha = 1$ (for a matrix-vector product) or 2 (for the Cholesky). Additionally, we estimate unknown parameters of the covariance function.

**Parameter estimation. Problem settings.** Let $n$ be the number of spatial measurements $Z$ located irregularly across a given geographical region at locations $s := \{s_1, \ldots, s_n\} \in \mathbb{R}^d$, $d \geq 1$. These data are frequently modeled as a realization from a stationary Gaussian spatial random field. Specifically, we let $Z = \{Z(s_1), \ldots, Z(s_n)\}^\top$, where $Z(s)$ is a Gaussian random field. Then, we assume that $Z$ has mean zero and a stationary parametric covariance function $C(h; \theta) = \text{cov}\{Z(s), Z(s+h)\}$, where $h \in \mathbb{R}^d$ is a spatial lag vector and $\theta \in \mathbb{R}^q$ is the unknown parameter vector of interest. To infer unknown parameters $\theta$ we maximize the Gaussian log-likelihood function:

$$L(\theta) = -\frac{n}{2} \log(2\pi) - \frac{1}{2} \log |C(\theta)| - \frac{1}{2} Z^\top C(\theta)^{-1} Z,$$

(1.1)

where $C(\theta)_{ij} = C(s_i - s_j; \theta)$, $i,j = 1, \ldots, n$. The maximum likelihood estimator of $\theta$ is the value $\hat{\theta}$ that maximizes (1.1). When the sample size $n$ is large, the evaluation of (1.1) becomes challenging, due to the computation of the quadratic form and log-determinant of the $n$-by-$n$ dense covariance matrix $C(\theta)$. Indeed, this requires $O(n^2)$ memory and $O(n^3)$ computational steps. Hence, scalable and efficient methods that can process larger sample sizes are needed.

**Definition 1.1.** An H-matrix approximation with the maximal rank $k$ of the exact log-likelihood $L(\theta)$ is defined by $\tilde{L}(\theta; k)$:

$$\tilde{L}(\theta; k) = -\frac{n}{2} \log(2\pi) - \sum_{i=1}^{n} \log \{\tilde{L}_{ii}(\theta; k)\} - \frac{1}{2} v(\theta)^\top v(\theta),$$

(1.2)

where $\tilde{L}(\theta; k)$ is an H-matrix approximation of the Cholesky factor $L(\theta)$ with the maximal rank $k$ in sub-blocks, $C(\theta) = L(\theta)L(\theta)^\top$, and $v(\theta)$ is the solution of the linear system $L(\theta; k)v(\theta) = Z$. 

The argument \( k \) in \( \tilde{L}(\theta; k) \) indicates that the fixed rank strategy, i.e., each sub-block in the covariance matrix has a rank equal to \( k \) or smaller, is used. In this case we cannot say anything about the accuracy in each sub-block (see more in Remark 2.2).

To maximize \( \tilde{L}(\theta; k) \) in (1.2), we use Brent-Dekker method \([15, 57]\). We note that maximization of the log-likelihood function is an ill-posed problem, since even very small perturbations in the covariance matrix \( C(\theta) \) may result in large perturbations in the log-determinant and in the log-likelihood. A possible remedy is to take a higher rank \( k \).

**Features of the \( \mathcal{H} \)-matrix approximation.** Other advantages for applying the \( \mathcal{H} \)-matrix technique are the following:

1. The class of \( \mathcal{H} \)-matrices is wide, for instance, it includes the classes of low-rank and sparse matrices;

2. \( C(\theta)^{-1}, C(\theta)^{1/2}, \det C(\theta), \) Cholesky decomposition, the Schur complement and, many others can be computed in the \( \mathcal{H} \)-matrix format \([29]\);

3. The \( \mathcal{H} \)-matrix technique is well-studied and has solid theory, many examples, and multiple sequential and parallel implementations, no specific MPI or OpenMP knowledge is needed;

4. The \( \mathcal{H} \)-matrix accuracy is controllable by the rank, \( k \). The full rank gives an exact representation;

5. The \( \mathcal{H} \)-matrix format preserves the structure (in contrast, for instance, to sparse matrices) after the Cholesky decomposition and the inverse have been computed, although with possibly a larger rank.

**Related work.** Stationary covariance functions, which have block Toeplitz or block circulant structure can be resolved with the Fast Fourier Transform (FFT) with the computing cost \( O(n \log n) \) \([23]\). However, this approach either does not work for data measured at irregularly spaced locations or requires expensive, non-trivial modifications.

Recently, a large amount of research has been devoted to approximating large covariance matrices: for example, covariance tapering \([22, 37, 65]\), likelihood approximations in both the spatial \([69]\) and spectral \([20]\) domains, latent processes such as Gaussian predictive processes \([6]\) and fixed-rank kriging \([17]\), and Gaussian Markov random-field approximations \([21, 44, 62, 63]\); see \([70]\) for a review. Each of these methods has its strengths and drawbacks \([67, 68, 71]\). Some other ideas include the nearest-neighbor Gaussian process models \([18]\), multiresolution Gaussian process models \([55]\), equivalent
kriging [40], multi-level restricted Gaussian maximum likelihood estimators [16], and hierarchical low-rank approximations [36]. Bayesian approaches to identify unknown or uncertain parameters could be also applied [49, 52, 53, 56, 59–61].

Previous results [29] show that the $H$-matrix technique is very stable when approximating the covariance matrix itself [4, 5, 12, 33, 38, 47, 64], its inverse [4, 10], its Cholesky decomposition [8, 9], and the conditional covariance matrix [29, 45, 46].

Recently, the maximum likelihood estimator for parameter-fitting Gaussian observations with a Matérn covariance matrix was computed via a framework for unstructured observations in two spatial dimensions, which allowed the evaluation of the log-likelihood and its gradient with computational complexity $O(n^{3/2})$, where $n$ was the number of observations; the method relied on the recursive skeletonization factorization procedure [34, 50]. However, the consequences of the approximation on the maximum likelihood estimator were not studied.

In [12] the authors computed the exact solution of Gaussian process regression via replacing the kernel matrix by a data-sparse approximation, called an $H^2$-matrix technique. An $H^2$-matrix approximation has $O(kn)$ computational complexity and storage cost, where $k$ is a parameter controlling the accuracy of the approximation.

Memory efficient kernel approximation was offered by [66]. The authors first make the observation that the structure of shift-invariant kernels changes from low-rank to block-diagonal (without any low-rank structure) when varying the scale parameter. Based on this observation, the authors propose a new kernel approximation algorithm – Memory Efficient Kernel Approximation, which considers both low-rank and clustering structure of the kernel matrix. They show that the resulting algorithm outperforms state-of-the-art low-rank kernel approximation methods regarding speed, approximation error, and memory usage.

In [5] the authors estimated the covariance matrix of a set of normally distributed random vectors. To overcome large numerical issues in high-dimensional regime, they computed (dense) inverses of sparse covariance matrices by the usage of $H$-matrices.

In [33] the authors offered methods for rapid computation of separable expansions for the approximation of random fields. In particular, they suggested the pivoted Cholesky method and provided a-posteriori error estimate in the trace norm. Low-rank tensor techniques in kriging are introduces in [54]. BigQUIC method for sparse inverse covariance estimation for a million variables was introduced in [35]. This method can solve 1 million dimensional $\ell_1$ regularized Gaussian MLE problems. In [64] the authors applied $H$-matrices to linear inverse problems in large-scale geostatistics. The $H^2$ matrix tech-
nique for solving large-scale stochastic linear inverse problems with applications in the subsurface modeling was demonstrated in [4]. From [3] we learned that typical covariance functions can be hierarchically factored into a product of block low-rank updates of the identity matrix, yielding an \(O(n \log n)\) algorithm for inversion. Additionally, the authors demonstrated that this factorization enables the evaluation of the determinant.

The structure of the paper. In Section 2 we introduce the methodology and algorithms. We review the \(H\)-matrix technique, the \(H\)-matrix approximations of Matérn covariance functions and Gaussian likelihood functions. In Section 3 we estimate the memory storage and computing costs. In Section 4 we describe software installation details, procedures of the HLIBCov code and the algorithm for parameter estimation. Estimation of unknown parameters is reported in Section 5. The best practices are listed in Section 6. We end the paper with a conclusion in Section 7. The auxiliary \(H\)-matrix details are provided in the Appendix A.

2 Methodology and algorithms

2.1 Matérn covariance functions

The class of Matérn covariance functions [27] is very widely used in spatial statistics, geostatistics, machine learning, image analysis, and other areas. The Matérn form of spatial correlations was introduced into statistics as a flexible parametric class [51], with one parameter determining the smoothness of the underlying spatial random field [32].

Let \(s\) and \(s'\) are any two spatial locations, and \(h := |s - s'|\). The Matérn class of covariance functions is defined as

\[
C(h; \theta) = \frac{\sigma^2}{2^{\nu-1} \Gamma(\nu)} \left( \frac{h}{\ell} \right)^\nu K_\nu \left( \frac{h}{\ell} \right),
\]

(2.1)

with \(\theta = (\sigma^2, \ell, \nu)^T\); where \(\sigma^2\) is the variance; \(\nu > 0\) controls the smoothness of the random field, with larger values of \(\nu\) corresponding to smoother fields; and \(\ell > 0\) is a spatial range parameter. Here \(K_\nu\) denotes the modified Bessel function of the second kind of order \(\nu\). When \(\nu = 1/2\), the Matérn covariance function reduces to the exponential covariance model and describes a rough field. The value \(\nu = \infty\) corresponds to a Gaussian covariance model that describes a very smooth, infinitely differentiable field. Random fields with a Matérn covariance function are \(\lfloor \nu - 1 \rfloor\) times mean square differentiable.

The Matérn covariance functions become simple [58] for \(\nu = p + 1/2\), where \(p\) is a non-negative integer. In this case the covariance function is a product of an exponential and a
polynomial of order $p$ \[1\]:

$$C_{y=3/2}(h) = \left( 1 + \frac{\sqrt{3}h}{\ell} \right) \exp\left( -\frac{\sqrt{3}h}{\ell} \right), \quad C_{y=5/2}(h) = \left( 1 + \frac{\sqrt{5}h}{\ell} + \frac{5h^2}{3\ell^2} \right) \exp\left( -\frac{\sqrt{5}h}{\ell} \right).$$

2.2 Hierarchical matrices

Hierarchical matrices have been described in detail [25, 28–31, 45]. Applications of the $H$-matrix technique to the covariance matrices can be found in [3, 4, 12, 33, 38, 39, 48]. Originally the $H$-matrix technique was developed for the approximation of stiffness matrices coming from partial differential and integral equations [14, 25, 28]. The idea behind $H$-matrices is to approximate blocks far from diagonal (i.e., far from the singularity) by low-rank matrices. The admissibility condition (criteria) is used to divide a given matrix into sub-blocks (Fig. 4) and define which sub-blocks can be approximated well by low-rank matrices and which not.

In general, covariance matrices are dense and, therefore, require $O(n^2)$ units of memory for the storage and $O(n^2)$ FLOPS for the matrix-vector multiplication. The $H$-matrix technique is defined as a hierarchical partitioning of a given matrix into sub-blocks (Fig. 4 right) followed by the further approximation of the majority of these sub-block by low-rank matrices (Fig. 3). After the decomposition of the matrix into sub-blocks has been
Figure 2: An example of $\mathcal{H}$-matrix $\tilde{C}$ with coarse sub-blocks (left) and fine (right). The adaptive rank arithmetic is used.

Figure 3: Examples of $\mathcal{H}$-matrix approximations of the Cholesky factor $\tilde{L}$. The adaptive rank arithmetic is used.
completed, an important question is how to compute the low-rank approximations. For this purpose the Adaptive Cross Approximation (or similar techniques like ACA+, HACA) algorithm [7,11,13,14,24] is used, which performs the approximations with a linear complexity $O(kn)$, where $n$ is the size of the sub-block.

**Remark 2.1.** The $H$-matrix approximation error may destroy the symmetry, therefore we do the following trick:

$$C := \frac{1}{2} (C + C^T).$$

**Definition 2.1.** We let $I$ be an index set. A hierarchical decomposition of $I$ we will call the cluster tree $T_I$ (see Fig. 4 and in Appendix). A hierarchical division of the index set product, $I \times I$, into sub-blocks (Fig. 4, right) is called block cluster tree $T_{I \times I}$. The set of $H$-matrices is

$$\mathcal{H}(T_{I \times I}; k) := \{ C \in \mathbb{R}^{I \times I} | \text{rank}(C|_b) \leq k \text{ for all admissible blocks } b \text{ of } T_{I \times I} \},$$

where $k$ is the maximum rank. Here, $C|_b = (c_{ij})_{(i,j) \in b}$ denotes the matrix block of $C = (c_{ij})_{i,j \in I}$ corresponding to sub-block $b \in T_{I \times I}$ (see Appendix).

Blocks that satisfy the admissibility condition can be approximated by low-rank matrices; see [28]. An $H$-matrix approximation of $C$ is denoted by $\tilde{C}$.  

Figure 4: (left) A cluster tree $T_I$ and (right) a block cluster tree $T_{I \times I}$ (matrix decomposition into sub-blocks).

**Remark 2.2.** (Fixed and adaptive rank strategies) For each sub-block we can either fix the accuracy $\varepsilon$ or the maximal approximative rank $k$. In the first case we speak about the adaptive rank strategy, i.e., where accuracy in the spectral norm for each sub-block is $\varepsilon$. In the second variant we speak about the fixed rank strategy, i.e., each sub-block has maximal rank $k$. The last variant does not provide us with the reliable estimate, but makes it easier to estimate the total computing cost and memory storage. We write $\mathcal{H}(T_{I \times I}; \varepsilon)$ and $\tilde{\mathcal{L}}(\theta; \varepsilon)$ for the adaptive ranks, and $\mathcal{H}(T_{I \times I}; k)$ and $\tilde{\mathcal{L}}(\theta; k)$ for the fixed. The fixed rank strategy is useful for a priori evaluations of the computational resources and storage memory. The
adaptive rank strategy is preferable for practical approximations and is useful when the accuracy in each sub-block is crucial.

In Fig. 5 (left) we demonstrate boxplots for different ranks $k \in \{3, 7, 9, 12\}$. In the experiment we run 100 replicates. One can see that boxplots for $k \in \{7, 9, 12\}$ are not changing, therefore there is no point to increase the rank, i.e. rank 7 is sufficient. Fig. 5 shows that the larger is the number of observations the better is estimation of the unknown parameter.

![Boxplot Diagram]

Figure 5: (left) Dependence of the boxplots for $\nu$ on the $H$-matrix rank, $n = 16,000$; (right) Convergence of boxplots for $\nu$ with increasing $n$, 100 replicates.

In Fig. 6 (left) we demonstrate the shape of the negative log-likelihood, of the quadratic form and of the log-determinant $\log(|C|)$. The true value of the parameter $\theta^* = \ell = 12$. In Fig. 6 (right) one can see three log-likelihood functions - the exact one and two others computed with different ranks 7 and 17. One can see, that all three plots look very similar.

2.3 Parallel hierarchical matrix technique

We used the parallel $H$-matrix library HLIBpro [26,41,43] to approximate the Matérn covariance matrix, to compute a Cholesky factorization, solve a linear system, calculate the determinant and a quadratic form. HLIBpro is fast, robust and efficient; see theoretical parallel complexity in Table 1. Here $|V(T)|$ denotes the number of vertices, $|L(T)|$ the number of leaves in the block-cluster tree $T = T_{1 \times 1}$, and $n_{\min}$ the size of a block when we stop further division into sub-blocks (see Section 2.2). Usually $n_{\min} = 32$ or 64, since a very deep hierarchy slows down computations.
Figure 6: (left) Shape of the log-likelihood function; (right) Three negative log-likelihood functions: the exact, and computed with $\mathcal{H}$-matrix ranks 7 and 17. One can see that even with rank 7 one can achieve very accurate results.

Table 1: Parallel complexity of the main linear operations in HLIBpro on $p$ cores.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Parallel Complexity $[41]$ (Shared Memory)</th>
</tr>
</thead>
<tbody>
<tr>
<td>build $\tilde{C}$</td>
<td>$\mathcal{O}(\log n) + \mathcal{O}(</td>
</tr>
<tr>
<td>store $\tilde{C}$</td>
<td>$\mathcal{O}(kn\log n)$</td>
</tr>
<tr>
<td>$\tilde{C} \cdot z$</td>
<td>$\mathcal{O}(k\log n) + \frac{n}{\sqrt{p}}$</td>
</tr>
<tr>
<td>$\alpha A \oplus \beta B$</td>
<td>$\mathcal{O}(kn\log n)$</td>
</tr>
<tr>
<td>$\alpha A \odot B \oplus \beta \tilde{C}$</td>
<td>$\mathcal{O}(n\log n) + \mathcal{O}(C_{\alpha\beta}(T)\setminus V(T))$</td>
</tr>
<tr>
<td>$\tilde{C}^{-1}$</td>
<td>$\mathcal{O}(n\log n) + \mathcal{O}(n n_{\text{min}}^2), n_{\text{min}} \approx 64$</td>
</tr>
<tr>
<td>$\mathcal{H}$-Cholesky $L$</td>
<td>$\mathcal{O}(n\log n) + \mathcal{O}(\frac{n^2\log n}{n^{1/d}}), d = 1,2,3$</td>
</tr>
<tr>
<td>determinant $</td>
<td>\tilde{C}</td>
</tr>
</tbody>
</table>

3 Memory storage, computing time and convergence

The Kullback-Leibler divergence (KLD) $D_{KL}(P\|Q)$ is a measure of information lost when distribution $Q$ is used to approximate $P$. For multivariate normal distributions $(\mu_0, \Sigma)$ and $(\mu_1, \Sigma)$ it is defined as follows

$$D_{KL}(\Sigma, \tilde{\Sigma}) = 0.5 \left( \text{tr}(\tilde{\Sigma}^{-1}\Sigma) + (\mu_1 - \mu_0)\top \tilde{\Sigma}^{-1}(\mu_1 - \mu_0) - n - \ln \left( \frac{|\Sigma|}{|\tilde{\Sigma}|} \right) \right)$$

In Tables 2 and 3 we show dependence of KLD and two matrix errors on the $\mathcal{H}$-matrix rank $k$ for Matérn covariance function with parameters $\ell = \{0.25,0.75\}$ and $\nu = \{0.5,1.5\}$,
computed on the domain $G = [0,1]^2$. All errors are under control, besides the last column. The ranks $k = 10, 12$ are not sufficient to approximate the inverse, the error $\|C(\tilde{C})^{-1} - I\|_2$ is large. A remedy is to take a larger rank.

<table>
<thead>
<tr>
<th>$k$</th>
<th>KLD $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
<th>$|C - \tilde{C}|_2$ $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
<th>$|C(\tilde{C})^{-1} - I|_2$ $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td>2.6e-3</td>
<td>0.2</td>
<td>7.7e-4</td>
<td>7.0e-4</td>
<td>6.0e-2</td>
<td>3.1</td>
</tr>
<tr>
<td>12</td>
<td>5.0e-4</td>
<td>2e-2</td>
<td>9.7e-5</td>
<td>5.6e-5</td>
<td>1.6e-2</td>
<td>0.5</td>
</tr>
<tr>
<td>15</td>
<td>1.0e-5</td>
<td>9e-4</td>
<td>2.0e-5</td>
<td>1.1e-5</td>
<td>8.0e-4</td>
<td>0.02</td>
</tr>
<tr>
<td>20</td>
<td>4.5e-7</td>
<td>4.8e-5</td>
<td>6.5e-7</td>
<td>2.8e-7</td>
<td>2.1e-5</td>
<td>1.2e-3</td>
</tr>
<tr>
<td>50</td>
<td>3.4e-13</td>
<td>5e-12</td>
<td>2.0e-13</td>
<td>2.4e-13</td>
<td>4e-11</td>
<td>2.7e-9</td>
</tr>
</tbody>
</table>

Table 2: Convergence of the $\mathcal{H}$-matrix approximation error vs the $\mathcal{H}$-matrix rank $k$, Matérn covariance with parameters $\ell = \{0.25, 0.75\}$ and $\nu = 0.5$, domain $G = [0,1]^2$, $\|C(\ell=0.25,0.75)\|_2 = \{212, 568\}$.

<table>
<thead>
<tr>
<th>$k$</th>
<th>KLD $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
<th>$|C - \tilde{C}|_2$ $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
<th>$|C(\tilde{C})^{-1} - I|_2$ $\ell = 0.25$</th>
<th>$\ell = 0.75$</th>
</tr>
</thead>
<tbody>
<tr>
<td>20</td>
<td>0.12</td>
<td>2.7</td>
<td>5.3e-7</td>
<td>2e-7</td>
<td>4.5</td>
<td>72</td>
</tr>
<tr>
<td>30</td>
<td>3.2e-5</td>
<td>0.4</td>
<td>1.3e-9</td>
<td>5e-10</td>
<td>4.8e-3</td>
<td>20</td>
</tr>
<tr>
<td>40</td>
<td>6.5e-8</td>
<td>1e-2</td>
<td>1.5e-11</td>
<td>8e-12</td>
<td>7.4e-6</td>
<td>0.5</td>
</tr>
<tr>
<td>50</td>
<td>8.3e-10</td>
<td>3e-3</td>
<td>2.0e-13</td>
<td>1.5e-13</td>
<td>1.5e-7</td>
<td>0.1</td>
</tr>
</tbody>
</table>

Table 3: Convergence of the $\mathcal{H}$-matrix approximation error vs the $\mathcal{H}$-matrix rank $k$, Matérn covariance with parameters $\ell = \{0.25, 0.75\}$ and $\nu = 1.5$, domain $G = [0,1]^2$, $\|C(\ell=0.25,0.75)\|_2 = \{720, 1068\}$.
Figure 7 shows that the $\mathcal{H}$-matrix storage cost is almost not changing for different parameters $\ell = \{0.15,...,2.2\}$ (on the left) and $\nu = \{0.3,...,1.3\}$ on the right. The computational domain is $[32.4,43.4] \times [-84.8,72.9]$ with $n=2,000$.

Figure 7: (left) Dependence of the matrix size on (left) the covariance length $\ell$; (right) the smoothness $\nu$ for two different accuracies in $\mathcal{H}$-matrix sub-blocks $\epsilon = \{1e^{-4},1e^{-6}\}$, $n=2,000$ locations in the domain $[32.4,43.4] \times [-84.8,72.9]$.

In Fig. 8 we plotted convergence of $\|C - \tilde{C}\|$ in the Frobenius and spectral norms vs. the rank $k$ for different smoothness parameters and covariance lengths. In Fig. 9 we plotted $\|C - \tilde{C}\|_2$ vs. rank $k$ for different $\nu$ and covariance lengths.
4 Software Installation and Numerical Examples

4.1 Introduction to HLIBpro

This Section contains shortened information taken from [www.hlib.com](http://www.hlib.com). The software library HLIBpro is a robust parallel C++ implementation of the $H$-matrix technique and $H$-matrix arithmetics, developed by Ronald Kriemann [41]. HLIBpro supports both the shared and distributed memory architectures, but in this work we use only the shared memory version. Intel’s Threading Building Blocks (TBB) are used for parallelisation on shared memory systems. HLIBpro is free for academic purposes, is distributed as a compiled code (no source code is available). Originally HLIBpro was developed for solving FEM and BEM problems [26, 43]. In this work we extend the applicability of HLIBpro to dense covariance matrices and log-likelihood functions.

4.2 Installation of HLIBpro

HLIBCov uses the functionality of HLIBpro, therefore, first, the reader needs to install HLIBpro. **Hardware.** All of the numerical experiments herein are performed on a Dell workstation with 20 processors (40 cores) and total 128 GB RAM.

**Operation system.** All the following steps are described for Linux systems (in our case Ubuntu 14.04). We did not test it for MacOS and for Windows operation systems. Since HLIBcov is just few c++ procedures (only one file loglikelihood.cc), the applicability of

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**Figure 9:** Convergence of the $H$-matrix approximation errors for $\nu = \{0.5, 1, 1.5\}$; (left) covariance length 0.1 and (right) covariance length 0.5.
HLIBCov is limited only by the applicability of HLIBpro. To install HLIBpro on MacOS and on Windows we refer to www.HLIBpro.com for further details.

**Software.** To use the binary distribution of HLIBpro, the following libraries have to be available together with HLIBpro: LAPACK/BLAS, Threading Building Blocks, Boost, zlib, FFTW3, Scotch, METIS, SCons (see details on www.HLIBpro.com). Additionally, the HLIBCov requires the GNU Scientific Library (GSL). We use GSL-1.16, the installation steps are described on https://www.gnu.org/software/gsl/. The GSL library provides maximization algorithms and Bessel functions. The reader can easily replace GSL on his own optimization library. The Bessel functions are also available in other packages.

To install LAPACK, zlib, FFTW3, Scotch and SCons using the standard package tools, you may run

```bash
$ aptitude install liblapack-dev zlib1g-dev libfftw3-dev libscotch-dev
$ aptitude install scons
```

A manual installation of TBB, Boost, Scotch and METIS requires appropriate modification of variables tbb_dir, boost_dir, scotch_dir andmetis_dir in the SConstruct file. After that, compile everything within the HLIBpro installation directory with

```bash
$ scons
$ export LD_LIBRARY_PATH=$LD_LIBRARY_PATH:<path_to_HLIBpro>/lib
```

The last terminal command is usually not necessary when using the provided SCons makefiles.

**An alternative - Using the Auxiliary Library Package.** As an alternative to the manual installing of all required packages (listed above), the HLIBpro provides a .tar.gz file with all auxiliary library packages. Download and unpack it into your HLIBpro directory:

```bash
$ cd HLIBpro-2.6
$ tar -xzvf HLIBpro-2.6-Linux-aux.tgz
```

Create a copy of SConstruct file, since the original SConstruct file will be overwritten. Afterwards, by calling scons you may compile the HLIBpro examples using the provided libraries.

**Adding HLIBCov to HLIBpro.** For default settings/paths, there is no need to change SConstruct file. To compile HLIBCov, add the following line to /examples/SConscript

```bash
$ examples.append(cxxenv.Program('loglikelihood.cc'))
$ cd ...
$ scons
$ cd examples/
$ ./loglikelihood
```
Input of the HLIBCov

The input contains in the first line the total number of locations $N$. Lines $2, \ldots, N+1$ contain the coordinates $x_i, y_i$, and the measurement value. An example:

<p>| | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>0.1</td>
<td>0.2</td>
</tr>
<tr>
<td>2</td>
<td>0.1</td>
<td>0.3</td>
</tr>
<tr>
<td>3</td>
<td>0.2</td>
<td>0.4</td>
</tr>
</tbody>
</table>

The HLIBpro requires neither a list of finite elements nor a list of edges. On the github we provide few input files examples of different size.

Output of the HLIBCov

The main output is the three identified parameter values $\theta = (\ell, \nu, \sigma^2)$. The auxiliary output may include a lot of details: $\mathcal{H}$-matrix details (the maximal rank $k$, the maximal accuracy in each sub-block, Frobenius and spectral norms of $\tilde{C}, L, L^{-1}, \|I - LL^{-1}\|_F$). Additionally, you can also print out iterations of the maximization algorithm. An example an output file below contains two iterations: index, $\nu$, $\ell$, $\sigma^2$, $\tilde{L}$, and the residual of the iterative method:

<p>| | | | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>0.27</td>
<td>2.4</td>
<td>1.30</td>
<td>$L = 1762.1$</td>
<td>$\text{TOL} = 0.007$</td>
</tr>
<tr>
<td>2</td>
<td>0.276</td>
<td>2.41</td>
<td>1.29</td>
<td>$L = 1757.2$</td>
<td>$\text{TOL} = 0.009$</td>
</tr>
</tbody>
</table>

If the iterative process is converging, then the last row will contain the solution $\theta^* = (\ell^*, \nu^*, \sigma^2)$. While computing error boxes, the output file will contain $M$ solutions $(n, \ell^*, \nu^*, \sigma^2)$, where $M$ is the number of replicants:

<p>| | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>4000</td>
<td>$5.4e-1$</td>
<td>$8.2e-2$</td>
</tr>
<tr>
<td>2</td>
<td>4000</td>
<td>$5.3e-1$</td>
<td>$8.3e-2$</td>
</tr>
<tr>
<td>3</td>
<td>4000</td>
<td>$5.5e-1$</td>
<td>$8.1e-2$</td>
</tr>
</tbody>
</table>

4.3 The preprocessing C++ function

Below we list the C++ code, which reads the data and initializes all required C++ objects (loglikelihood.cc).

```
1 vector< double * > vertices;
2 TScalarVector rhs;
3 INIT();
4 in >> N;
5 cout << " reading " << N << " coordinates" << endl;
6 vertices.resize( N );
7 rhs.set_size( N );
8 double x, y, v=0.0;
9 for ( int i = 0; i < N; ++i ){
10   in >> x >> y >> v;
11   vertices[i] = new double[ dim ];
12   vertices[i][0] = x; vertices[i][1] = y;
```

The C++ code computing the maximum of the log-likelihood function (loglikelihood.cc).

```cpp
double call_compute_max_likelihood(TScalarVector Z, double nu, double covlength, double sigma2, TBlockClusterTree* bct, TClusterTree* ct, TVector< double >* vertices, double output[3])
{
    gsl_function F;
    int status; iter = 0, max_iter = 200; smy_f_params params;
    FILE* f1; double size;
    const gsl_multimin_fminimizer_type *T = gsl_multimin_fminimizer_nmsimplex2;
    gsl_multimin_fminimizer *s = NULL; gsl_vector *ss, *x;
    gsl_multimin_function minex_func;
    params.bct = bct; params.ct = ct; params.Z = Z; params.nu = nu;
    params.covlength = covlength; params.sigma2 = sigma2; params.vertices = vertices;
    /* Starting point */
    x = gsl_vector_alloc(3); gsl_vector_set (x, 0, nu);
    gsl_vector_set (x, 1, covlength); gsl_vector_set (x, 2, sigma2);
    /* Set initial step sizes to 0.1 */
    ss = gsl_vector_alloc (3);
    gsl_vector_set (ss, 0, 0.02); // for nu
    gsl_vector_set (ss, 1, 0.04); // for theta
    gsl_vector_set (ss, 2, 0.01); // for sigma2
    /* Initialize method and iterate */
    minex_func.n = 3; // dimension
    minex_func.params = &params;
    s = gsl_multimin_fminimizer_alloc (T, 3); /* or 2? like in the example */
    gsl_multimin_fminimizer_set (s, &minex_func, x, ss);
    do{
        status = gsl_multimin_fminimizer_iterate(s);
        if (status) break;
        size = gsl_multimin_fminimizer_size (s); // for stopping criteria
        status = gsl_multimin_fminimizer_status(s);
        if (status == GSL_SUCCESS) printf("converged to minimum at ");
    } while (status == GSL_CONTINUE && iter < max_iter);
    output[0]= gsl_vector_get(s->x, 0); // nu
    output[1]= gsl_vector_get(s->x, 1); // theta
    output[2]= gsl_vector_get(s->x, 2); // sigma2
    return status;
}
```

Below we list the C++ code, which computes the value of the log-likelihood for given parameters (loglikelihood.cc).

```cpp
double eval_logli (const gsl_vector *sol, void* p)
{
    pmy_f_params params;
    double nu = gsl_vector_get(sol, 0);
    double length = gsl_vector_get(sol, 1);
    double sigma2 = gsl_vector_get(sol, 2);
    unique_ptr< TProgressBar > progress( verbose(2) ? new TConsoleProgressBar : nullptr );
    params = (pmy_f_params)p;
    double vec[3];
    CALL_EVAL_MAX_LIKELIHOOD(nu, covlength, sigma2, bct.get(), ct.get(), vertices, vec);
    return vec[0];
}
```
TScalarVector rhs = (params->Z);
TBlockClusterTree* bct = (params->bct); TClusterTree* ct = (params->ct);
vector< double > > vertices= (params->vertices);
double err2=0.0, nugget = 1.0e-4, s = 0.0;
auto acc = fixed_prec(1e-5); int dim = 2, N = 0;
TCovCoeffFn coefffn(length,nu,sigma2,nugget,vertices,ct->perm_i2e(),ct->perm_i2e());
TACAPlus< real_t > aca( & coefffn );
TACAPlus< real_t > real_t > h_builder( & coefffn, & aca );
// enable coarsening during construction
h_builder.set_coarsening( false );
auto & A = b_builder.build( bct, acc, progress.get() );
N=A->cols();
auto & A_copy = A->copy();
auto options = fac_options_t( progress.get() );
options.eval = point_wise; // Extreme important
auto & A_inv = ldl_inv( A_copy.get(), acc, options );
for ( int i = 0; i < N; ++i ) {
const auto v = A_copy->entry( i, i );
  s = s + log(v); // for
}

HLIBpro License
To receive the license file HLIBpro.lic send an email with your user login name and the machine name, on which you plan to run the code, to Ronald Kriemann rok@mis.mpg.de and request a license file. The machine name can be received with the terminal command hostname. The received license file must be placed either in the directory where the final program will be executed, e.g. in the examples/ subdirectory, or in the directory pointed to by the environment variable HLIBpro_LIC_PATH. You can add this variable to the system paths

```
export HLIBpro_LIC_PATH=<path_to_HLIBpro.lic>
```

5 Numerical experiments with synthetic data
We performed numerical experiments with simulated data to recover the true values of the parameters of the Matérn covariance matrix, known to be \((\ell^*,\nu^*,\sigma^*) = (1.0,0.5,1.0)\).

5.1 Generation of the synthetic data
To build \(M\) various data sets (\(M\) replicates) with \(n \in \{128,64,\ldots,4,2\} \times 1000\) locations, we generate a large vector \(Z_0\) with \(n_0 = 2 \cdot 10^6\) locations and randomly sample \(n\) points from it. Note, that if locations are very close to each other, then the covariance matrix could be
singular or it will be difficult to compute the Cholesky factorization.

To generate the random data \( Z_0 \in \mathbb{R}^{n_0} \) we compute the \( \mathcal{H} \)-Cholesky factorization of \( C(1.0,0.5,1.0) = \tilde{L} \tilde{L}^\top \). Then we evaluate \( Z_0 = \tilde{L} \xi \), where \( \xi \in \mathbb{R}^{n_0} \) is a normal vector with zero mean and unit variance. We generate \( Z_0 \) only once. After that we run our optimization algorithm and try to identify (recover) the “unknown” parameters \((\ell^*, \nu^*, \sigma^2)\). The resulting boxplots for \( \ell \) and \( \nu^2 \) over \( M = 100 \) replicates are illustrated in Fig. 10. One can see that the variance (or uncertainty) is decreasing with increasing of \( n \). The green line indicates the true values.

![Boxplots for \( \ell \) and \( \nu \) over \( M = 100 \) replicates.](image)

Figure 10: Synthetic data with known parameters \((\ell^*, \nu^*, \sigma^2) = (0.5, 1, 0.5)\). Boxplots for \( \ell \) and \( \nu \) for \( n = 1,000 \times \{64, 32, \ldots, 4, 2\} \), 100 replicates.

To identify all three parameters simultaneously a three-dimensional optimization problem is solved. The maximal number of iterations is set to 200, the residual to \( 10^{-6} \). The behavior and accuracy of the boxplots depend on the hierarchical matrix rank, the maximum number of iterations to achieve a certain threshold, the threshold (or residual) itself, the initial guess, the step size in each parameter in the maximization algorithm, and the maximization algorithm. All replicates of \( Z \) are sampled from the same generated vector of size \( n_0 = 2 \cdot 10^6 \).

In Table 4 we demonstrate the almost linear storage cost (columns 3 and 6) and computing time (columns 2 and 5).

To get an idea about the shape of the log-likelihood function, and about its components we draw Fig. 11. This picture helps us to understand the behavior of the iterative
Table 4: Computing time and storage cost for parallel $\mathcal{H}$-matrix approximation; number of cores is 40, $\hat{\nu} = 0.33$, $\hat{\ell} = 0.65$, $\hat{\sigma}^2 = 1.0$. $\mathcal{H}$-matrix accuracy in each sub-block for both $\tilde{C}$ and $\tilde{L}$ is $10^{-5}$.

<table>
<thead>
<tr>
<th>$n$</th>
<th>comp. time</th>
<th>size MB</th>
<th>kB/dof</th>
<th>comp. time</th>
<th>size MB</th>
<th>$| I - (\tilde{L}\tilde{L}^T)^{-1}\tilde{C}|_2$</th>
</tr>
</thead>
<tbody>
<tr>
<td>32,000</td>
<td>3.3</td>
<td>162</td>
<td>5.1</td>
<td>2.4</td>
<td>172.7</td>
<td>$2.4 \cdot 10^{-3}$</td>
</tr>
<tr>
<td>128,000</td>
<td>13.3</td>
<td>776</td>
<td>6.1</td>
<td>13.9</td>
<td>881.2</td>
<td>$1.1 \cdot 10^{-2}$</td>
</tr>
<tr>
<td>512,000</td>
<td>52.8</td>
<td>3420</td>
<td>6.7</td>
<td>77.6</td>
<td>4150</td>
<td>$3.5 \cdot 10^{-2}$</td>
</tr>
<tr>
<td>2,000,000</td>
<td>229</td>
<td>14790</td>
<td>7.4</td>
<td>473</td>
<td>18970</td>
<td>$1.4 \cdot 10^{-1}$</td>
</tr>
</tbody>
</table>

Figure 11: Dependence of the negative log-likelihood and its ingredients on parameter $\ell$, parameter $\nu = 0.5$ is fixed, the true value of covariance length is 0.2337. Exponential covariance, $n = 66049$, rank $k = 16$, $\sigma^2 = 1$.

optimization method, contribution of the log-determinant, of the quadratic functional. One can see that the log-likelihood is almost flat and it may be necessary a lot of iteration steps to find the minimum.

**Nugget** $\tau^2$. Assume $|\tilde{C}| \neq 0$. It is known \[19\], that for a small perturbation matrix $E$ it is hold

$$ \frac{\| (\tilde{C} + E)^{-1} - (\tilde{C})^{-1} \|}{\| \tilde{C}^{-1} \|} \leq \kappa(\tilde{C}) \cdot \frac{\| E \|}{\| \tilde{C} \|} = \frac{\kappa(\tilde{C}) \tau^2}{\| \tilde{C} \|}, $$

where $\kappa(\tilde{C})$ is the condition number of $\tilde{C}$, and $E = \tau^2 I$. Or, substituting $\kappa(\tilde{C}) := \| \tilde{C}^{-1} \| \cdot \| \tilde{C}^{-1} \|$, obtain

$$ \frac{\| (\tilde{C} + \tau^2 I)^{-1} - (\tilde{C})^{-1} \|}{\| \tilde{C}^{-1} \|} \leq \tau^2 \| \tilde{C}^{-1} \|. $$
Table 5: Comparison of three log-likelihood functions computed with three different $H$-matrix accuracies \{$10^{-7}, 10^{-9}, 10^{-11}$\}. Exponential covariance function discretized in the domain $[32.4, 43.4] \times [-84.8, -72.9]$, $n=32,000$ locations. Columns corresponds to different covariance lengths $\{0.001, \ldots, 0.1\}$.

<table>
<thead>
<tr>
<th>$\ell$</th>
<th>0.001</th>
<th>0.005</th>
<th>0.01</th>
<th>0.02</th>
<th>0.03</th>
<th>0.05</th>
<th>0.07</th>
<th>0.1</th>
</tr>
</thead>
<tbody>
<tr>
<td>$-\tilde{L}(\ell;10^{-7})$</td>
<td>44657</td>
<td>36157</td>
<td>36427</td>
<td>40522</td>
<td>45398</td>
<td>68450</td>
<td>70467</td>
<td>90649</td>
</tr>
<tr>
<td>$-\tilde{L}(\ell;10^{-9})$</td>
<td>44585</td>
<td>36352</td>
<td>36113</td>
<td>41748</td>
<td>47443</td>
<td>60286</td>
<td>70688</td>
<td>90615</td>
</tr>
<tr>
<td>$-\tilde{L}(\ell;10^{-11})$</td>
<td>44529</td>
<td>37655</td>
<td>36390</td>
<td>42020</td>
<td>47954</td>
<td>60371</td>
<td>72785</td>
<td>90639</td>
</tr>
</tbody>
</table>

From the last equation one can see that if $a \leq \|\hat{C}^{-1}\| \leq b$, where $a$, $b$ are some positive constants, then the relative error will be of order $\tau^2$. Figure 12 (left) demonstrates three negative log-likelihood functions. One is computed with the nugget 0.01, another with 0.005 and the third with 0.001. As one can see, for this particular example, the behavior of likelihood is preserved, and the minimum is not changing (or changing very slightly). Figure 12 (right) is just a zoomed version of the left picture.

Figure 12: (left) Dependence of the log-likelihood on parameter $\ell$ with nuggets ($\{0.01, 0.005, 0.001\}$) for Gaussian covariance. (right) Zoom of the left figure near minimum; $n=2000$ random points from moisture example, rank $k=14$, $\sigma^2=1$. 
6 Best practices

In this section we list our recommendations and warnings.

1. Initialize variables and the license with `INIT();`

2. For practical computations use the adaptive rank arithmetics.

3. For the input it is enough to define a file with 3 columns: location coordinates \((x, y)\) and the observed value, no triangles, no edges are required;

4. If two locations are coincide or very close to each other, then the matrix will be close to singular or singular. As a result, it will be hard to compute the Cholesky factorization. A remedy is to improve the quality of locations.

5. Construction of \(H\)-matrices permutes the indices of the degrees of freedom. Therefore, there are two permutation mappings: `ct->perm.e2i()` and `ct->perm.i2e()`.
For example, to prepare the observation vector \(Z\), computed on a grid, for multiplication with the hierarchical matrix \(L^{-1}\) (or \(C^{-1}\)) use `ct->perm.e2i()->permute(&Z)`
To find to which grid node corresponds the \(i\)-th row (column), use `ct->perm.i2e()`

6. HLIBpro uses smart pointers, e.g., `std::unique_ptr`. If the user leaves a code block those smart pointers will automatically delete the object they manage. This simplifies programming and makes it much more secure.

7. By default, the \(H\)-Cholesky or \(H-LU\) factorizations are always DAG based. You can turn this off by setting `HLIB::CFG::Arith::use_dag = false;`

8. By default, HLIBpro uses all available computing cores. To perform computations on 16 cores, use `HLIB::CFG::set_nthreads(16)` at the beginning of the program (after command `INIT()`).

9. Since HLIBpro is working for 2D and 3D, there are only very minor changes in HLIBCov to move from 2D locations to 3D. Replace `dim=2` to `dim=3` in `TCoordinate coord(vertices, dim);`
Then add `>> z` to `in >> x >> y >> z >> v;`

10. Finalize all computations with `DONE();`;
7 Conclusion

We extended functionality of the parallel $H$-matrix library HLIBpro for application in spatial statistics and in parameters inference. This new extension allows us to work with large covariance matrices. The first novelty is that we approximated the joint multivariate Gaussian likelihood function and have found its maxima in the $H$-matrix format. This maxima estimated unknown parameters ($\ell$, $\nu$, and $\sigma^2$) of the covariance model. The second novelty is that the new code is parallel, highly efficient and written in C++ language. With the $H$-matrix technique we reduced the storage cost and the computing cost (Tables 3, 4) of the log-likelihood function dramatically, from cubic to almost linear. We demonstrated a synthetic example, where we were able to identify the true parameters of the covariance model. With the suggested method we were able to compute the log-likelihood function for 2,000,000 locations just in a few minutes on a powerful desktop machine (Table 4). The $H$-matrix technique allows us to increase the spatial resolution, to handle more measurements, consider larger regions and identify more parameters simultaneously.

Acknowledgments

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References


[18] A. Datta, S. Banerjee, A. O. Finley, and A. E. Gelfand. Hierarchical nearest-neighbor Gauss-
sian process models for large geostatistical datasets. *Journal of the American Statistical Association*, 0(0a):00–00, 2015.


A (Block) Cluster trees and admissibility condition

Let $I$ be an index set of all degrees of freedom. Denote for each index $i \in I$ corresponding to a basis function $b_i$ the support $G_i := \text{supp} b_i \subset \mathbb{R}^d$. Now we define two trees which are necessary for the definition of hierarchical matrices. These trees are labeled trees where the label of a vertex $t$ is denoted by $\hat{t}$.

Definition A.1. (Cluster Tree $T_{I \times I}$)

A finite tree $T_I$ is a cluster tree over the index set $I$ if the following conditions hold:

- $I$ is the root of $T_I$ and a subset $\hat{t} \subseteq I$ holds for all $t \in T_I$.
- If $t \in T_I$ is not a leaf, then the set of sons $\text{sons}(t)$ contains disjoint subsets of $I$ and the subset $\hat{t}$ is the disjoint union of its sons, $\hat{t} = \bigcup_{s \in \text{sons}(t)} \hat{s}$.
- If $t \in T_I$ is a leaf, then $|\hat{t}| \leq n_{\text{min}}$ for a fixed number $n_{\text{min}}$.

Definition A.2. (Block Cluster Tree $T_{I \times I}$) Let $T_I$ be a cluster tree over the index set $I$. A finite tree $T_{I \times I}$ is a block cluster tree based on $T_I$ if the following conditions hold:

- $\text{root}(T_{I \times I}) = I \times I$.
- Each vertex $b$ of $T_{I \times I}$ has the form $b = (\tau, \sigma)$ with clusters $\tau, \sigma \in T_I$.
- For each vertex $(\tau, \sigma)$ with $\text{sons}(\tau, \sigma) \neq \emptyset$, we have
  
  $\text{sons}(\tau, \sigma) = \begin{cases} 
  (\tau, \sigma') : \sigma' \in \text{sons}(\sigma), & \text{if } \text{sons}(\tau) = \emptyset \land \text{sons}(\sigma) \neq \emptyset \\
  (\tau', \sigma) : \tau' \in \text{sons}(\tau), & \text{if } \text{sons}(\tau) \neq \emptyset \land \text{sons}(\sigma) = \emptyset \\
  (\tau', \sigma') : \tau' \in \text{sons}(\tau), \sigma' \in \text{sons}(\sigma), & \text{otherwise}
  \end{cases}$
• The label of a vertex \((\tau, \sigma)\) is given by \(\widehat{(\tau, \sigma)} = \tau \times \sigma \subseteq I \times I\).

We can see that \(\text{root}(T_{I \times I}) = I \times I\). This implies that the set of leaves of \(T_{I \times I}\) is a partition of \(I \times I\).

## B Rank-\(k\) Adaptive Cross Approximation (ACA)

An \(\mathcal{H}\)-matrix contains many sub-blocks, which can be well approximated by low-rank matrices. How to compute these low-rank approximations? The truncated singular value decomposition is accurate, but very slow. HLIBpro uses the Adaptive Cross Approximation method (ACA) \[24\] and its improved modifications such as ACA+ and HACA \[7, 11, 13\].

**Definition B.1.** Let \(R \in \mathbb{R}^{p \times q}\), then factorization

\[
R = AB^\top, \quad \text{where } A \in \mathbb{R}^{p \times k}, \quad B \in \mathbb{R}^{q \times k}, \quad k \ll \min\{p, q\} \in \mathbb{N}.
\] (B.1)

we will call a low-rank representation.

Note that any matrix of rank \(k\) can be represented in the form (B.1).

Suppose that \(b\) is a sub-block in the block-cluster tree, and \(R := C|_b\). Suppose it is known that \(R\) could be approximated by a rank-\(k\) matrix. We explain below how to compute \(R\) in the form (B.1). ACA is especially effective for assembling low-rank matrices. It requires to compute only \(k\) columns and \(k\) rows of the matrix under consideration and, thus, has the computational cost \(k(p + q)\). In \[24\] it is proved that if there exists a sufficiently good low-rank approximation, then there also exists a cross approximation with almost the same accuracy in the sense of the spectral norm. The ACA algorithm computes vectors \(a_\ell\) and \(b_\ell\) which form \(\hat{R} = \sum_{\ell=1}^{k} a_\ell b_\ell^\top\) such that \(\|R - \hat{R}\| \leq \varepsilon\), where \(\varepsilon\) is the desired accuracy \[11\] \[14\]. In \[13\], the reader can also find different counterexamples when the standard ACA algorithm does not work. Here we present the standard version of the ACA algorithm.

**Algorithm B.1.** ACA algorithm

```plaintext
begin
\text{/* input is a required accuracy } \varepsilon \text{ and a function to compute } R_{ij} */;
\text{/* output is matrix } \hat{R} */;
k = 0; \hat{R} = 0;
S = \emptyset; T = \emptyset; \text{/* sets of row and column indices */}
```
do
  Take a row \( i^* \notin S \);
  Subtract \( R_{i^*j} := R_{i^*j} - \hat{R}_{i^*j} \), \( j = 1..q \);
  Find \( \max_j |a_{i^*j}| \neq 0, j < q \). Suppose it lies in column \( j^* \);
  Compute all elements \( b_{ij^*} \) in column \( j^* \), \( i < p \);
  Subtract \( R_{ij^*} := R_{ij^*} - \hat{R}_{ij^*} \), \( i = 1..p \);
  \( k := k + 1 \); \( S := S \cup \{i^*\} \); \( T := T \cup \{j^*\} \);
  Compute \( \tilde{R} = \hat{R} + a_{i^*} \cdot b_{j^*}^\top \); \( / \ast \) it is rank \( k \) approximation\( / \ast \)
  if(\( \|a_{i^*} \cdot b_{j^*}^\top \|_2 \leq \varepsilon \cdot \|a_1 \cdot b_1^\top \|_2 \)) return \( \tilde{R} \);
  Find \( \max_i |b_{ij^*}|, i < p, i \neq i^* \). The row where it lies is a new row \( i^* \);
  \( \text{until}(k < k_{\text{max}}) \)
  return \( \tilde{R} \);
end;

Note that the algorithm does not compute the whole matrix \( R \). The subtraction is
done only from the elements under consideration, i.e. row \( a_\ell \) and column \( b_\ell, \ell = 1,...,k \).

Remark B.1. Further optimisation of the ACA algorithm can be done by the truncated
SVD. Suppose that a factorisation of matrix \( R = AB^\top \), \( A \in \mathbb{R}^{p \times K}, B \in \mathbb{R}^{q \times K} \), is found by
ACA. Suppose also that the rank of \( R \) is \( k, k < K \). Then one can apply the truncated SVD
algorithm to compute \( R = U \Sigma V^\top \) requiring \( O((p+q)K^2 + K^3) \) FLOPS.